

Delta Modeling in Practice

A Fredhopper Case Study

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VaMoS 2012



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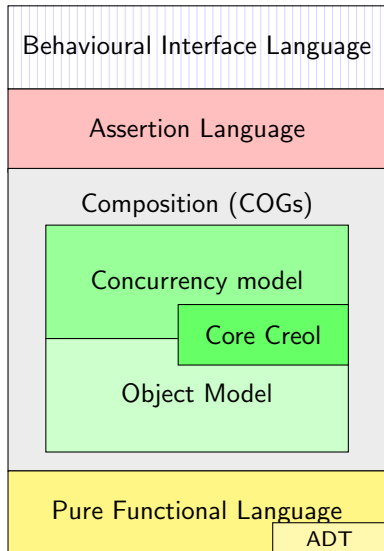
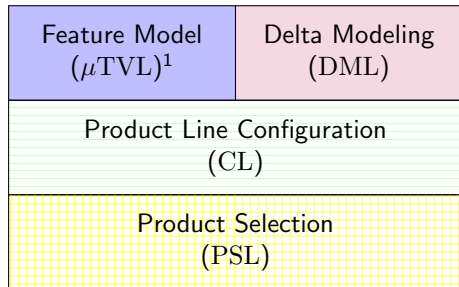
Delta Modeling Workflow (DMW)

- ▶ DMW is a method for modeling software product lines (SPL)
- ▶ Abstract Behavioral Specification (ABS) is an executable modeling language that implements delta modeling as its core paradigm for modeling adaptable systems
- ▶ We report on the development of an industrial scale SPL in ABS following DMW to guide the application of delta modeling in practice

HATS: Highly Adaptable & Trustworthy Software Using Formal Models

- ▶ FP7 project
- ▶ Develop a tool-supported formal software engineering methodology for the design, analysis and implementation of highly adaptable and trustworthy software systems
- ▶ Started 1 March 2009, 48 months runtime
- ▶ Integrated Project, academically driven
- ▶ 9 academic partners, 2 industrial research, 1 SME
- ▶ 8 countries

ABS Language Layers



¹Based on: A. Classen, Q. Boucher, P. Heymans. A Text-based Approach to Feature Modelling: Syntax and Semantics of TVL. SCP 2010.

Fredhopper

- ▶ eCommerce company based in Amsterdam,
- ▶ Flagship product – Fredhopper Access Server (FAS)
- ▶ FAS provides search and merchandising IT services for e-Commerce companies.
- ▶ FAS is a concurrent distributed object based system
- ▶ Clients include: Toys “R” Us, Clarks, ASOS, Debenhams, Philips

FAS Concepts

Data a set of product items to search from

Data Manager transforms customer's data into FAS input **data**

Query Engine receives query (e.g. "adidas shoes") and returns a subset of items satisfying that query.

Configuration how to present that subset of items (with add-ons like upsell promotions etc.)

Business Manager provides an interface for changing **configuration**

FAS Concepts

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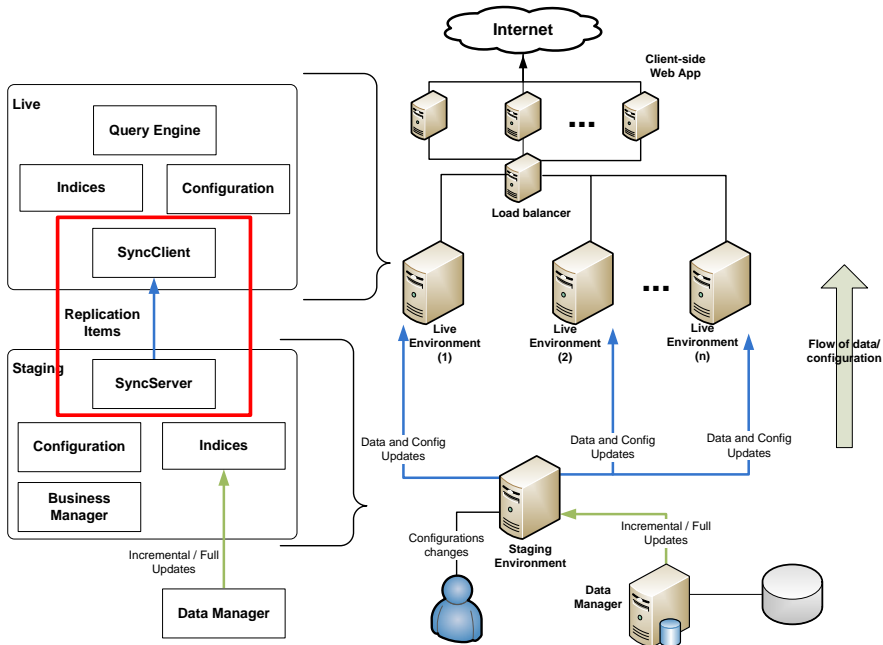
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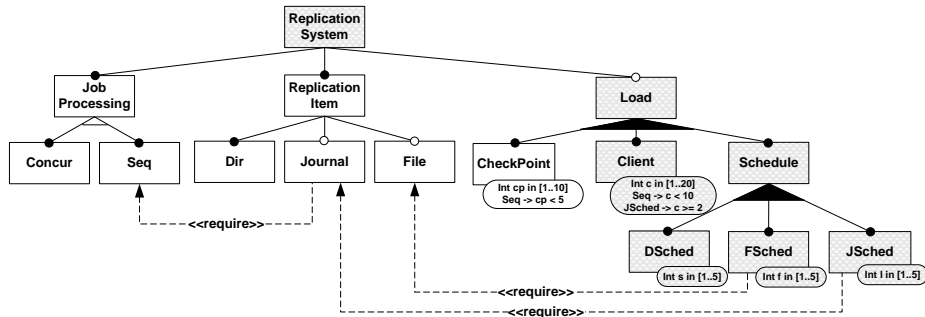
Replication System **updates** data and configuration, and **maintains** their consistency throughout a FAS deployment (multi-threaded)



Replication System Product Line

- ▶ synchronises **configurations** and **data** from the staging environment to all live environments
- ▶ consists of a server (SyncServer) and one or more clients (SyncClient)
- ▶ desired properties: deadlock and livelock freedom, replication consistency

Feature diagram of the Replication System



Variability (Features)

Replication Item data dependent variation

Job Processing resource dependent variation

Load resource dependent variation

```
root RS {
  group allof {
    JobProcessing { ... },
    ReplicationItem { ... },
    Load {
      group [1..3] {
        CheckPoint { ... },
        Client { Int c in [1 .. 20]; Seq -> c < 10; JSched -> c >= 2; },
        Schedule {
          group [1..3] {
            DSched { Int s in [1..5]; },
            FSched { Int f in [1..5]; require: File; },
            JSched { Int l in [1..5]; require: Journal; }
          }
        }
      }
    }
  }
}
```

[◀ Configuration](#)[◀ Feature Interaction](#)[◀ Implement Feature Interaction](#)

Feature Load

Client Number of SyncClients

Schedule Number of locations in the file system at which changes to various replication item types are monitored: *DSched* for directories, *FSched* for file sets, and *JSched* for transaction journals.

Feature Interaction

Client and *JSched* require some extra implementation effort to make them interact properly:

$$\exists p : p \models \{Client\} \wedge p \models \{JSched\} \wedge p \not\models \{Client, JSched\}$$

▶ μ TVL model

Structural Categorization

Load and *Schedule* are subcategorizations with no semantics:

$$\forall p : p \models F \iff p \models F \cup \{Load, Schedule\}$$

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SPL Specification

Consider only features *RS*, *Load*, *Client*, *Schedule*, *DSched*, *FSched* and *JSched*: SPL specification (Ψ', \models') where $\models' \subset \models$ and,

$$\Psi' = (\{\{RS\}, \emptyset, \{(RS, Load), (Load, Client), (Load, Schedule), (Schedule, DSched), (Schedule, FSched), (Schedule, JSched)\}, \emptyset, \emptyset)$$

Implement features in topological order of the feature diagram:

$\langle RS, Load, Client, Schedule, DSched, FSched, JSched \rangle$

Given SPL: $RS = (c, \emptyset, \emptyset, \emptyset)$, where c is a minimum core ABS model

```
class RSMMain { } { new RSMMain(); }
```

1 Implement feature *RS* with delta *RD*

```
delta RD {  
  adds data SchedType = Dir | File | Journal; adds type CId = Int; ...  
  adds class System(...) { ... }  
  modifies class RSMain {  
    adds Map<CP,Map<Fld,Content>> datas = map[...];  
    adds Map<SchedType,List<Schedule>> ss = map[...];  
    adds Set<CId> cids = set[...];  
    adds Unit run() {  
      Map<CP,Map<Fld,Content>> is = this.getDatas();  
      List<Schedule> ss = this.getSchedules(); Set<CId> cs = this.getCids();  
      new System(is,ss,cs); }  
    adds List<Schedule> getSchedules() { return lookup(ss,Dir); }  
    adds Map<CP,Map<Fld,Content>> getDatas() { return datas; }  
    adds Set<CId> getCids() { return cids; } } }
```

Step 1 results in SPL: $RS = (c, \{RD\}, \emptyset, \{(RD, \Phi')\})$:

- ② No need to implement feature *Load*
- ③ Implement feature *Client* with delta CD

```
delta CD(Int c) {  
  modifies class RSMMain {  
    modifies Set<CId> getCids() { return takeSet(RD.original(),c); }  
  }  
}
```

- 4 Implement feature *DSched* with delta DD

```
delta DD(Int s) {  
  modifies class RSMMain {  
    modifies List<Schedule> getSchedules() {  
      return take(lookup(RD.original(),Dir),s);  
    }  
  }  
}
```

- 5 Implement feature *FSched* with delta FD

```
delta FD(Int f) {  
  modifies class RSMMain {  
    modifies List<Schedule> getSchedules() {  
      return take(lookup(RD.original(),File),f);  
    }  
  }  
}
```

Conflict between FD and delta DD: resolution with delta DFD

```
delta DFD {  
  modifies class RSMMain {  
    modifies List<Schedule> getSchedules() {  
      return concatenate(DD.original(),FD.original());  
    }  
  }  
}
```

6 Implement feature *JSched* with delta JD

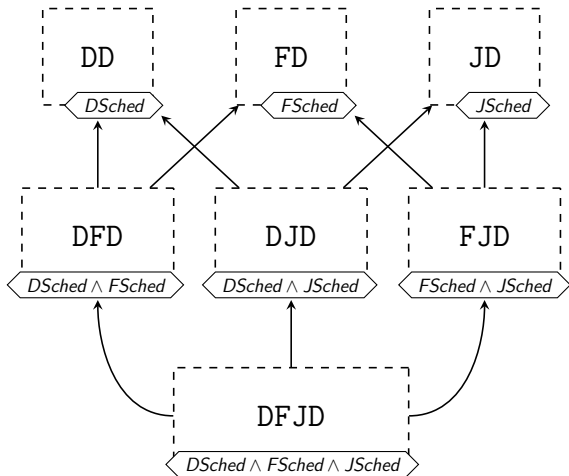
```
delta JD(Int I) {  
  modifies class RSMMain {  
    modifies List<Schedule> getSchedules() {  
      return take(lookup(RD.original(),Journal),I);  
    }  
  }  
}
```

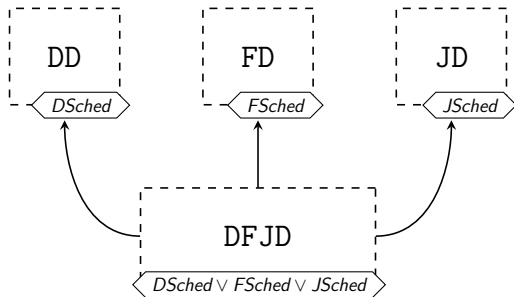
Implement feature interaction between *Client* and *JSched* with delta JCD

```
delta JCD {  
  modifies class RSMMain {  
    modifies Set<CId> getCids() {  
      Set<CId> cs = CD.original();  
      if (size(cs) == 1) {  
        cs = Insert(failSafe(),c);  
      }  
      return cs;  
    }  
  }  
}
```

- 7 Resolve conflict between *DSched*, *FSched* and *JSched* ...

Delta Modeling





- 7 Resolve conflict between *DSched*, *FSched* and *JSched* with delta DFJD

```
delta DFJD(Bool ds, Bool fs, Bool js) {  
  modifies class RSMMain {  
    modifies List<Schedule> getSchedules() {  
      List<Schedule> ss = Nil;  
      if (ds) { ss = DD.original(); }  
      if (fs) { ss = concatenate(ss,FD.original()); }  
      if (js) { ss = concatenate(ss,JD.original()); }  
      return ss;  
    }  
  }  
}
```


Product Line Configuration & Selection

```
productline RS {  
  features RS, Client, DSched, FSched, JSched;  
  delta RD when RS;  
  delta CD (Client.c) when Client after RD;  
  delta DD(DSched.s) when DSched after RD;  
  delta FD(FSched.f) when FSched after RD;  
  delta JD(JSched.l) when JSched after RD;  
  delta JCD when Client and JSched after CD, JD;  
  delta DFJD(DSched, FSched, JSched)  
    when DSched or FSched or JSched  
    after DD, FD, JD;  
}
```

```
product SingleClient (RS, Client{c=1}, DSched{s=3});  
product TwoClients (RS, Client{c=2}, DSched{s=3});
```

▶ μ TVL model

Complete Case Study

- ▶ \approx 5000 lines of codes
- ▶ 40 Classes
- ▶ 43 Interfaces
- ▶ 89 user-defined functions
- ▶ 17 user-defined data types
- ▶ 15 Features
- ▶ 10 (7) Deltas
- ▶ 96 (12024) Products

DMW

Completeness unambiguity, feature interaction, conflict resolution, error reduction

Flexibility three-way conflict resolution

Evolution relaxed empty initial product

Collaboration atomicity, noninterference

Tool support convert μTVL to Ψ

ABS

- ▶ Targeted **original**
- ▶ Feature Boolean
- ▶ Application condition*

So far...

- ▶ First account of using delta modeling to implement a system of industrial scale and of practical use
- ▶ Served as a test bed for DMW, providing feedback that was used to refine the workflow
- ▶ Evaluation of the practical applicability of ABS modeling language

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- ▶ First account of using delta modeling to implement a system of industrial scale and of practical use
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Upcoming...

Timing constraint modelling timed behavior

Resource constraint modelling underlying resources

Verification specifying and verifying replication consistency